### A LLT-like test for proving the primality of Fermat numbers.

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In 1876, Édouard Lucas discovered a method for proving that a number is prime or composite without searching its factors. His method was based on the properties of the Lucas Sequences. He first used his method for Mersenne numbers and proved that  $2^{127} - 1$  is a prime. In 1930, Derrick Lehmer provided a complete and clean proof. This test of primality for Mersenne numbers is now known as: Lucas-Lehmer Test (LLT).

Few people know that Lucas also used his method for proving that a Fermat number is prime or composite, still with an unclear proof. He used his method for proving that  $2^{2^6} + 1$  is composite. Lehmer did not provide a proof of Lucas' method for Fermat numbers.

This paper provides a proof of a LLT-like test for Fermat numbers, based on the properties of Lucas Sequences and based on the method of Lehmer. The seed (the starting value  $S_0$  of the  $\{S_i\}$  sequence) used here is 5, though Lucas used 6. In 1960, Kustaa Inkeri provided a full proof with seed 8.

Primality tests for special numbers are classified into N-1 and N+1categories, meaning that the numbers N-1 or N+1 can be completely or partially factored. Since many books talk about the LLT only in the N+1chapter for Mersenne numbers  $N = 2^q - 1$ , it seemed useful to remind that the LLT can also be used for numbers N such that N-1 is easy to factor, like Fermat numbers  $N = 2^{2^n} + 1$ , by providing a proof *àla* Lehmer.

#### Theorem 1

 $F_n=2^{2^n}+1 \ (n\geqslant 1)$  is a prime if and only if it divides  $S_{2^n-2}$ , where  $S_0=5$  and  $S_i=S_{i-1}^2-2$  for  $i=1,2,3,\ldots 2^n-2$ .

The proof is based on chapters 4 (The Lucas Functions) and 8.4 (The Lehmer Functions) of the book "Edouard Lucas and Primality Testing" of H. C. Williams (A Wiley-Interscience publication, 1998).

Chapter 1 explains how the (P, Q) parameters have been found. Then Chapter 2 provides the Lehmer theorems used for the proof. Then Chapter 3 and 4 provide the proof for:  $F_n$  prime  $\implies F_n \mid S_{2^n-2}$  and the converse, proving theorem 1. Chapter 5 provides numerical examples. The appendix in Chapter 6 provides first values of  $U_n$  and  $V_n$  plus some properties.

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## 1 Lucas Sequence with $P = \sqrt{R}$

Let  $S_0 = 5$  and  $S_i = S_{i-1}^2 - 2$ .  $S_1 = 23, S_2 = 527 = 17 \times 31, \dots$ 

It has been checked that: 
$$\begin{cases} S_{2^n-2} \equiv 0 \pmod{F_n} & \text{for } n = 1...4\\ S_{2^n-2} \neq 0 \pmod{F_n} & \text{for } n = 5...14 \end{cases}$$

Here after, we search a Lucas Sequence  $(U_m)_{m \ge 0}$  and its companion  $(V_m)_{m \ge 0}$ with (P, Q) that fit with the values of the  $S_i$  sequence.

We define the Lucas Sequence  $V_m$  such that:

$$V_{2k+1} = S_k \tag{1}$$

Thus we have: 
$$\begin{cases} V_2 = S_0 = 5\\ V_4 = S_1 = 23\\ V_8 = S_2 = 527 \end{cases}$$

If (4.2.7) page 74 ( $V_{2n} = V_n^2 - 2Q^n$ ) applies, we have:  $\begin{cases}
V_4 = V_2^2 - 2Q^2 \\
V_8 = V_4^2 - 2Q^4
\end{cases}$ 

and thus:  $Q = \sqrt[2]{\frac{V_2^2 - V_4}{2}} = \sqrt[4]{\frac{V_4^2 - V_8}{2}} = \pm 1$ .

With (4.1.3) page 70 (  $V_{n+1} = PV_n - QV_{n-1}$  ), and with:

$$\begin{cases} V_0 = 2 \\ V_1 = P \\ V_2 = PV_1 - QV_0 = P^2 - 2Q \end{cases}$$

we have:  $P = \sqrt{V_2 + 2Q} = \sqrt{7}$  or  $\sqrt{3}$ . In the following we consider:  $(P,Q) = (\sqrt{7},1)$ .

As explained by Williams page 196, "all of the identity relations [Lucas functions] given in (4.2) continue to hold, as these are true quite without regard as to whether P, Q are integers".

So, like Lehmer, we define  $P = \sqrt{R}$  such that R = 7 and Q = 1 are coprime integers and we define (Property (8.4.1) page 196):

$$\overline{V}_n = \begin{cases} V_n & \text{when } 2 \mid n \\ V_n / \sqrt{R} & \text{when } 2 \nmid n \end{cases} \quad \overline{U}_n = \begin{cases} U_n / \sqrt{R} & \text{when } 2 \mid n \\ U_n & \text{when } 2 \nmid n \end{cases}$$

in such a way that  $\overline{V}_n$  and  $\overline{U}_n$  are always integers.

Tables 1 to 5 give values of  $U_i$ ,  $V_i$ ,  $\overline{U}_i \pmod{F_n}$ ,  $\overline{V}_i \pmod{F_n}$ , with  $(P,Q) = (\sqrt{7},1)$ , for n = 1, 2, 3, 4.

### 2 Lehmer theorems

Like Lehmer, let define the symbols (where (a/b) is the Legendre symbol):

$$\begin{cases} \varepsilon = \varepsilon(p) = (D/p) \\ \sigma = \sigma(p) = (R/p) \\ \tau = \tau(p) = (Q/p) \end{cases}$$

The 2 following formulas (from page 77) will help proving properties:

$$(4.2.28) \quad 2^{m-1}U_{mn} = \sum_{i=0}^{\lfloor \frac{m}{2} \rfloor} {m \choose 2i+1} D^{i}U_{n}^{2i+1}V_{n}^{m-(2i+1)}$$
$$(4.2.29) \quad 2^{m-1}V_{mn} = \sum_{i=0}^{\lfloor \frac{m}{2} \rfloor} {m \choose 2i} D^{i}U_{n}^{2i}V_{n}^{m-2i}$$

Property (8.4.2) page 196 :

If p is an odd prime and 
$$p \nmid Q$$
, then: 
$$\begin{cases} \overline{U}_p \equiv (D/p) \pmod{p} \\ \overline{V}_p \equiv (R/p) \pmod{p} \end{cases}$$

#### **Proof:**

Since p is a prime, and by Fermat little theorem, we have:  $2^{p-1} \equiv 1 \pmod{p}$ .

• By (4.2.28), with m = p and n = 1, since  $U_1 = 1$  and  $V_1 = P$ , we have:

$$2^{p-1}U_p = \sum_{i=0}^{\frac{p-1}{2}} {p \choose 2i+1} D^i U_1^{2i+1} V_1^{p-(2i+1)}$$
$$2^{p-1}U_p = {p \choose 1} P^{p-1} + {p \choose 3} DP^{p-3} + \dots + {p \choose p} D^{\frac{p-1}{2}} P^0$$

Since  $\binom{p}{i} \equiv 0 \pmod{p}$  when 0 < i < p and  $\binom{p}{p} = 1$ , we have:

$$U_p = \overline{U}_p \equiv D^{\frac{p-1}{2}} \equiv (D_p) \pmod{p}$$

• By (4.2.29), with m = p and n = 1, since  $U_1 = 1$  and  $V_1 = P$ , we have:

$$2^{p-1}V_p = \sum_{i=0}^{\frac{p-1}{2}} {p \choose 2i} D^i U_1^{2i} V_1^{p-2i}$$

$$2^{p-1}V_p = \binom{p}{0}P^p + \binom{p}{2}DP^{p-2} + \dots + \binom{p}{p-1}D^{\frac{p-1}{2}}P$$

Since  $\binom{p}{0} = 1$ , and  $\binom{p}{i} \equiv 0 \pmod{p}$  when 0 < i < p, we have:

$$V_p \equiv P^p \text{ and } \overline{V}_p \equiv P^{p-1} \equiv R^{\frac{p-1}{2}} \equiv \left(\frac{R}{p}\right) \pmod{p}$$

Property (8.4.3) page 197 :

 $p \text{ odd prime and } p \nmid Q \Longrightarrow p \mid \overline{U}_{p-\sigma\varepsilon}$ 

Proof

By (4.2.28) with  $n = 1, V_1 = P$ , since p is a prime and (R, Q) = 1, we have:

• With: m = p + 1

$$2^{p}U_{p+1} = \sum_{i=0}^{\frac{p+1}{2}} {p+1 \choose 2i+1} D^{i}P^{p-2i}$$

$$2^{p}U_{p+1} = \binom{p+1}{1}P^{p} + \binom{p+1}{3}DP^{p-2} + \dots + \binom{p+1}{p}D^{\frac{p-1}{2}}P + \binom{p+1}{p+2}D^{\frac{p+1}{2}}P^{-1}$$

$$2^{p}U_{p+1} = (p+1)P^{p} + (p+1)p[\dots] + (p+1)D^{\frac{p-1}{2}}P + 0D^{\frac{p+1}{2}}P^{-1}$$

$$2^{p}U_{p+1} = P^{p} + D^{\frac{p-1}{2}}P + p[\dots] = P[(P^{2})^{\frac{p-1}{2}} + D^{\frac{p-1}{2}}] + p[\dots]$$

$$\frac{2^{p}U_{p+1}}{P} = 2^{p}\overline{U}_{p+1} \equiv R^{\frac{p-1}{2}} + D^{\frac{p-1}{2}} \equiv (R_{p}) + (D_{p}) = \sigma(p) + \varepsilon(p) \pmod{p}$$

Thus, if  $\sigma\varepsilon=\sigma(p)\times\varepsilon(p)=-1$  , then  $p\mid\overline{U}_{p+1}=\overline{U}_{p-\sigma\epsilon}$  .

• With: m = p - 1 :

$$2^{p-2}U_{p-1} = \sum_{i=0}^{\frac{p-1}{2}} {p-1 \choose 2i+1} D^i P^{p-2(i+1)}$$

$$\begin{split} 2^{p-2}U_{p-1} &= \binom{p-1}{1}P^{p-2} + \binom{p-1}{3}DP^{p-4} + \ldots + \binom{p-1}{p-2}D^{\frac{p-3}{2}}P + \binom{p-1}{p}D^{\frac{p-1}{2}}P^{-1} \\ 2^{p-2}U_{p-1} &= (p-1)P^{p-2} + (p-1)DP^{p-4} + \ldots + (p-1)D^{\frac{p-3}{2}}P + 0D^{\frac{p-1}{2}}P^{-1} \\ \frac{2^{p-2}U_{p-1}}{P} &\equiv -[P^{p-3} + DP^{p-5} + \ldots + D^{\frac{p-3}{2}}] \equiv -\frac{P^{p-1} - D^{\frac{p-1}{2}}}{P^2 - D} \pmod{p} \\ 2^{p-2}\overline{U}_{p-1}(P^2 - D) &\equiv -(P^2)^{\frac{p-1}{2}} + D^{\frac{p-1}{2}} \equiv \varepsilon(p) - \sigma(p) \pmod{p} \\ \text{Thus, if } \sigma\varepsilon &= \sigma(p) \times \varepsilon(p) = 1 \text{, then } p \mid \overline{U}_{p-1} = \overline{U}_{p-\sigma\epsilon} \text{.} \end{split}$$

#### Property (8.4.4) page 197

If p is an odd prime and  $p \nmid Q$ , then:  $V_{p-\sigma\varepsilon} \equiv 2\sigma Q^{\frac{1-\sigma\varepsilon}{2}} \pmod{p}$ .

**Theorem 2 (8.4.1)** If p is an odd prime and  $p \nmid QRD$ , then:

$$\begin{cases} p \mid \overline{V}_{\frac{p-\sigma\epsilon}{2}} & when \quad \sigma = -\tau \\ p \mid \overline{U}_{\frac{p-\sigma\epsilon}{2}} & when \quad \sigma = \tau \end{cases}$$

**Definition (8.4.2) page 197 of**  $\omega(m)$  : For a given m, denote by  $\omega = \omega(m)$  the value of the least positive integer k such that  $m \mid \overline{U}_k$ . If  $\omega(m)$  exists,  $\omega(m)$  is called the **rank of apparition** of m.

Theorem 3 (8.4.3)

$$\begin{cases} If k \mid n, then \overline{U}_k \mid \overline{U}_n \ . \\ If m \mid \overline{U}_n, then \omega(m) \mid n \ . \end{cases}$$

**Theorem 4 (8.4.5)** If (m, Q) = 1, then  $\omega(m)$  exists.

**Theorem 5 (8.4.6)** If (N, 2QRD) = 1 and  $N \pm 1$  is the rank of apparition of N, then N is a prime.

**Theorem 6 (8.4.7)** If (N, 2QRD) = 1,  $\overline{U}_{N\pm 1} \equiv 0 \pmod{N}$  and  $\overline{U}_{\frac{N\pm 1}{q}} \neq 0 \pmod{N}$  for each distinct prime divisor q of  $N \pm 1$ , then N is a prime.

#### **Proof:**

Let  $\omega = \omega(N)$ . We see that  $\omega \mid N \pm 1$ , but  $\omega \nmid (N \pm 1)/q$ . Thus if  $q^{\alpha} \parallel N \pm 1$ , then  $q^{\alpha} \mid \omega$ . It follows that  $\omega = N \pm 1$  and N is a prime by Theorem 5 (8.4.6).

**3** 
$$F_n$$
 prime  $\Longrightarrow F_n \mid \overline{V}_{\frac{F_n-1}{2}}$  and  $F_n \mid S_{2^n-2}$ 

Let  $N=F_n=2^{2^n}+1$  with  $n\geq 1$  be an odd prime. Let:  $P=\sqrt{R}$  , R=7 , Q=1 , and  $D=P^2-4Q=3$  .

Hereafter we compute (3/N) and (7/N):

• 
$$(^{3}/_{N})$$
:  
Since: 
$$\begin{cases} N \text{ odd prime} \\ N = (4)^{2^{n-1}} + 1 \equiv 2 \pmod{3} \\ (N/_{3}) = (^{2}/_{3}) = -1 \\ (3/_{N}) = (N/_{3}) \times (-1)^{\frac{3-1}{2}\frac{N-1}{2}} \end{cases}$$
 then:  $(^{3}/_{N}) = -1$ .

•  $\left(\frac{7}{N}\right)$ : We have:  $\begin{cases} 2^3 \equiv 1 \pmod{7} \\ 2^{3a+b} \equiv 2^b \pmod{7} \end{cases}$ 

With  $2^n \equiv b \pmod{3}$ , we have:  $2^{2^n} + 1 \equiv 2^b + 1 \pmod{7}$ . Then we study the exponents of 2, modulo 3. We have:  $2^2 \equiv 1 \pmod{3}$ , and:

If 
$$n = 2m$$

$$\begin{cases}
2^{2m} \equiv 1 \pmod{3} \\
N = 2^{2^{2m}} + 1 \equiv 2^{1} + 1 \equiv 3 \pmod{7} \\
(N/7) = (3/7) = -1
\end{cases}$$
If  $n = 2m + 1$ 

$$\begin{cases}
2^{2m+1} \equiv 2 \pmod{3} \\
N = 2^{2^{2m+1}} + 1 \equiv 2^{2} + 1 \equiv 5 \pmod{7} \\
(N/7) = (5/7) = -1
\end{cases}$$

Finally, we have:  $(7_N) = (N_7)(-1)^{\frac{7-1}{2}2^{2^n}} = (N_7) = -1$ .

So we have: 
$$\begin{cases} \varepsilon = (D_{N}) = (3_{N}) = -1 \\ \sigma = (R_{N}) = (7_{N}) = -1 \\ \tau = (Q_{N}) = (1_{N}) = +1 \end{cases}$$

Since  $\sigma = -\tau$  ,  $\sigma \epsilon = +1$  , and  $F_n \nmid QRD$  with  $n \ge 1$ , then by Theorem 2 (8.4.1) we have:

$$F_n \text{ prime } \Longrightarrow F_n \mid \overline{V}_{\frac{F_n-1}{2}} = V_{2^{2^n-1}}$$

By (1) we have:  $V_{2^{k-1}} = S_{k-2}$  and thus, with  $k = 2^n$ :  $F_n \mid S_{2^n-2}$ .

# 4 $F_n \mid S_{2^n-2} \implies F_n$ is a prime

Let  $N = F_n$  with  $n \ge 1$ . By (1) we have:  $N \mid S_{2^n-2} \Longrightarrow N \mid V_{2^{2^n-1}}$ . And thus, by (4.2.6) page 74 ( $U_{2a} = U_a V_a$ ), we have:  $N \mid \overline{U}_{2^{2^n}}$ .

By (4.3.6) page 85: (  $(V_n,U_n)\mid 2Q^n$  for any n ), and since Q=1, then:  $(V_{2^{2^n}-1},\overline{U}_{2^{2^n}-1})=2$  and thus:  $N\nmid\overline{U}_{2^{2^n}-1}$  since N odd.

With  $\omega = \omega(N)$ , by Theorem 3 (8.4.3) we have :  $\omega \mid 2^{2^n}$  and  $\omega \nmid 2^{2^{n-1}}$ . This implies:  $\omega = 2^{2^n} = N - 1$ . Then N - 1 is the rank of apparition of N, and thus by Theorem 5 (8.4.6) N is a prime.

This test of primality for Fermat numbers has been communicated to the community of number theorists working on this area on mersenneforum.org (http://www.mersenneforum.org/showthread.php?t=2130) in May 2004, and the proof was finalized in September 2004.

Then, in a private communication, Robert Gerbicz provided a proof of the same theorem based on  $Q[\sqrt{21}]$ .

## 5 Numerical Examples

 $(\text{mod } F_2) \ S_0 = 5 \stackrel{1}{\mapsto} 6 \stackrel{2}{\mapsto} S_{2^2 - 2} \equiv 0 \\ (\text{mod } F_3) \ S_0 = 5 \stackrel{1}{\mapsto} 23 \stackrel{2}{\mapsto} 13 \stackrel{3}{\mapsto} 167 \stackrel{4}{\mapsto} 131 \stackrel{5}{\mapsto} 197 = -60 \stackrel{6}{\mapsto} S_{2^3 - 2} \equiv 0 \\ (\text{mod } F_4) \ S_0 = 5 \stackrel{1}{\mapsto} 23 \stackrel{2}{\mapsto} 527 \stackrel{3}{\mapsto} 15579 \stackrel{4}{\mapsto} 21728 \stackrel{5}{\mapsto} 42971 \stackrel{6}{\mapsto} 1864 \stackrel{7}{\mapsto} 1033 \stackrel{8}{\mapsto} 18495 \stackrel{9}{\mapsto} 27420 \stackrel{10}{\mapsto} 15934 \stackrel{11}{\mapsto} 2016 \stackrel{12}{\mapsto} 960 \stackrel{13}{\mapsto} 4080 \stackrel{14}{\mapsto} S_{2^4 - 2} \equiv 0$ 

## 6 Appendix: Table of $U_i$ and $V_i$

With n = 2, 3, 4, we have the following (not proven) properties (modulo  $F_n$ ):

$\overline{U}_{F_n-5}$	$\equiv$	5		$\int \overline{V}_{F_n-5}$	$\equiv$	-23
$\overline{U}_{Fn-4}$	$\equiv$	6		$\overline{V}_{F_n-4}$	$\equiv$	-4
$\overline{U}_{Fn-3}$	$\equiv$	1		$\overline{V}_{Fn-3}$	$\equiv$	-5
$\overline{U}_{Fn-2}$	$\equiv$	1		$\overline{V}_{Fn-2}$	$\equiv$	-1
$\overline{U}_{F_n-1}$	$\equiv$	0	•	$\overline{V}_{F_{n-1}}$	$\equiv$	-2
$\overline{U}_{Fn}$	$\equiv$	-1		$\overline{V}_{Fn}$	$\equiv$	-1
$\overline{U}_{F_n+1}$	$\equiv$	-1		$\overline{V}_{F_{n+1}}$	$\equiv$	-5
$\overline{U}_{F_n+2}$	$\equiv$	-6		$\overline{V}_{F_n+2}$	$\equiv$	-4
$\overline{U}_{F_n+3}$	≡	-5		$\overline{V}_{F_n+3}$	≡	-23

The values of  $\overline{U'}_n$  and  $\overline{V'}_n$   $(n \ge 1)$  with  $(P, Q) = (\sqrt{3}, -1)$  can be built by:

$$\begin{cases} \overline{U'}_{2n} = \overline{U}_{2n} \\ \overline{U'}_{2n+1} = \overline{V}_{2n+1} \end{cases} \begin{cases} \overline{V'}_{2n} = \overline{V}_{2n} \\ \overline{V'}_{2n+1} = \overline{U}_{2n+1} \end{cases}$$

Values of  $U_i$  and  $V_i$  in previous tables can be computed easily by the following PARI/gp programs:

 $U_{2j+1}$ : U0=1;U1=6; for(i=1,N, U0=5\*U1-U0; U1=5\*U0-U1; print(4\*i+1," ",U0); print(4\*i+1," ",U1))

i	$U_i$		$V_i$	
0	0	$\times \sqrt{7}$	2	
1	1		1	$\times \sqrt{7}$
2	1	$\times \sqrt{7}$	5	
3	6		4	$\times \sqrt{7}$
4	5	$\times \sqrt{7}$	23	
5	29		19	$\times \sqrt{7}$
6	24	$\times \sqrt{7}$	110	
7	139		91	$\times \sqrt{7}$
8	115	$\times \sqrt{7}$	527	
9	666		436	$\times \sqrt{7}$
10	551	$\times \sqrt{7}$	2525	
11	3191		2089	$\times \sqrt{7}$
12	2640	$\times \sqrt{7}$	12098	
13	15289		10009	$\times \sqrt{7}$
14	12649	$\times \sqrt{7}$	57965	
15	73254		47956	$\times \sqrt{7}$
16	60605	$\times \sqrt{7}$	277727	
17	350981		229771	$\times \sqrt{7}$
18	290376	$\times \sqrt{7}$	1330670	
19	1681651		1100899	$\times \sqrt{7}$
20	1391275	$\times \sqrt{7}$	6375623	
21	8057274		5274724	$\times \sqrt{7}$
22	6665999	$\times \sqrt{7}$	30547445	
23	38604719		25272721	$\times \sqrt{7}$
24	31938720	$\times \sqrt{7}$	146361602	
25	184966321	_	121088881	$\times \sqrt{7}$
26	153027601	$\times \sqrt{7}$	701260565	
27	886226886		580171684	$\times \sqrt{7}$
28	733199285	$\times \sqrt{7}$	3359941223	_
29	4246168109	_	2779769539	$\times \sqrt{7}$
30	3512968824	$\times \sqrt{7}$	16098445550	_
31	20344613659	_	13318676011	$\times \sqrt{7}$
32	16831644835	$\times \sqrt{7}$	77132286527	
33	97476900186	_	63813610516	$\times \sqrt{7}$
34	80645255351	$\times \sqrt{7}$	369562987085	
35	467039887271		305749376569	$\times \sqrt{7}$
36	386394631920	$\times \sqrt{7}$	1770682648898	
37	2237722536169		1464933272329	$\times \sqrt{7}$
38	1851327904249	$\times \sqrt{7}$	8483850257405	
39	10721572793574		7018916985076	$\times \sqrt{7}$
40	8870244889325	$\times \sqrt{7}$	40648568638127	

Table 1:  $P = \sqrt{7}$  , Q = 1

i	$\overline{U}_i \pmod{F_1}$	$\overline{V}_i \pmod{F_1}$
0	0	2
1	1	1
2	1	0
3	1	4
4	0	3
5	4	4
6	4	0
7	4	1
8	0	2

Table 2:  $P = \sqrt{7}$  , Q = 1 , Modulo  $F_1$ 

i	$\overline{U}_i \pmod{F_2}$	$\overline{V}_i \pmod{F_2}$
0	0	2
1	1	1
2	1	5
3	6	4
4	5	6
5	12	2
6	7	8
7	3	6
8	13	0
9	3	11
10	7	9
11	12	15
12	5	11
13	6	-4
14	1	-5
15	1	-1
16	0	-2
17	-1	-1
18	-1	-5
19	-6	-4
20	-5	11
21	5	15
22	10	9
23	14	11
24	4	0

Table 3:  $P=\sqrt{7}~,~Q=1$ , Modulo $F_2$ 

i	$\overline{U}_i \pmod{F_3}$	$\overline{V}_i \pmod{F_3}$
0	0	2
1	1	1
2	1	5
3	6	4
4	5	23
8	115	13
16	210	167
32	118	131
64	38	197
128	33	0
192	38	60
224	118	126
240	210	90
248	115	-13
252	5	-23
253	6	-4
254	1	-5
255	1	-1
256	0	-2
257	-1	-1
258	-1	-5
259	-6	-4
260	-5	-23

Table 4:  $P=\sqrt{7}~,~Q=1$ , Modulo $F_3$ 

i	$\overline{U}_i \pmod{F_4}$	$\overline{V}_i \pmod{F_4}$
2048	9933	15934
4096	567	2016
8192	28943	960
16384	63129	4080
32768	5910	0
65532	5	-23
65533	6	-4
65534	1	-5
65535	1	-1
65536	0	-2
65537	-1	-1
65538	-1	-5
65539	-6	-4
65540	-5	-23

Table 5:  $P=\sqrt{7}~,~Q=1$ , Modulo $F_4$